

# Performance and Energy Consumption Analysis of a Delay-Tolerant Network for Censorship-Resistant Communication

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June 25, 2015

# Outline

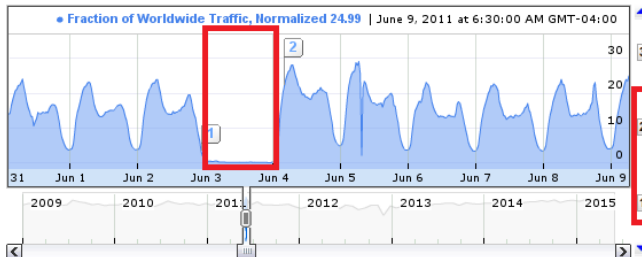
- 1 Introduction
- 2 1am Deployment
- 3 Performance and Overhead
- 4 Conclusion

# Surveillance and Censorship Threaten Privacy and Freedom of Speech on the Internet

Many countries are conducting active censorship at this moment...

# Surveillance and Censorship Threaten Privacy and Freedom of Speech on the Internet

Fraction of Worldwide Traffic, Normalized



3. Syria disconnects itself from the Internet for 40 minutes – ZDNet [\[Read More\]](#)

2012-7-19

2. Syrian government unplugs Internet for much of country – NBCnews.com [\[Read More\]](#)

2011-6-4



1. Syrian government unplugs Internet for much of country – NBCnews.com [\[Read More\]](#)

2011-6-2

# Surveillance and Censorship Threaten Privacy and Freedom of Speech on the Internet

TOP SECRET//SI//ORCON//NOFORN

Gmail facebook Hotmail Google Skype paltalk.com YouTube AOL mail

 (TS//SI//NF) **PRISM Collection Details** 

Current Providers

What Will You Receive in Collection (Surveillance and Stored Comms)?  
It varies by provider. In general:

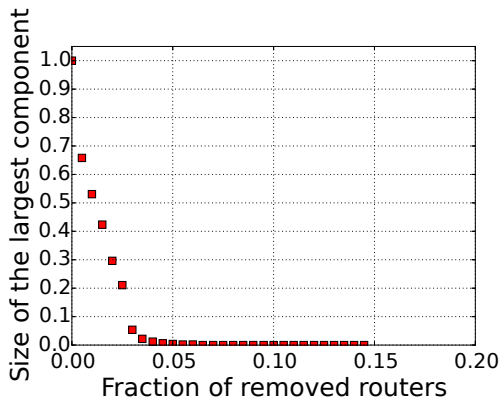
- Microsoft (Hotmail, etc.)
- Google
- Yahoo!
- Facebook
- PalTalk
- YouTube
- Skype
- AOL
- Apple

- E-mail
- Chat – video, voice
- Videos
- Photos
- Stored data
- VoIP
- File transfers
- Video Conferencing
- Notifications of target activity – logins, etc.
- Online Social Networking details
- **Special Requests**

Complete list and details on PRISM web page:  
Go PRISMFAA

TOP SECRET//SI//ORCON//NOFORN

# The Hierarchical Internet

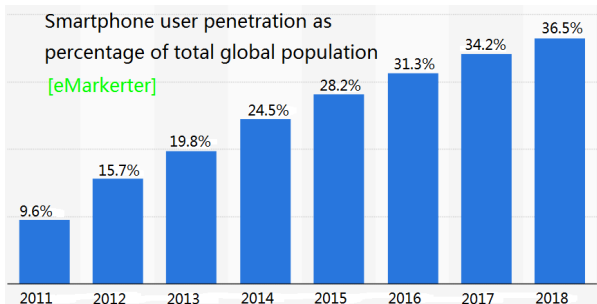


Fragmentation of the Internet as its most connected routers are removed. [\[Albert'00\]](#)

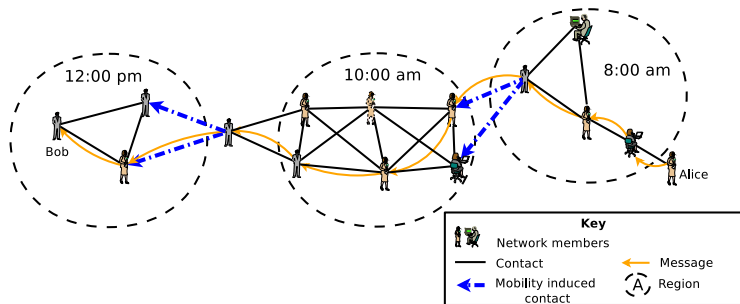
# Infrastructureless Networks Composed of Commodity Mobile Devices

## Advantages

- Very low technical/economical barrier to participation
- High availability



# Infrastructureless Networks: DTNs Composed of Commodity Mobile Devices

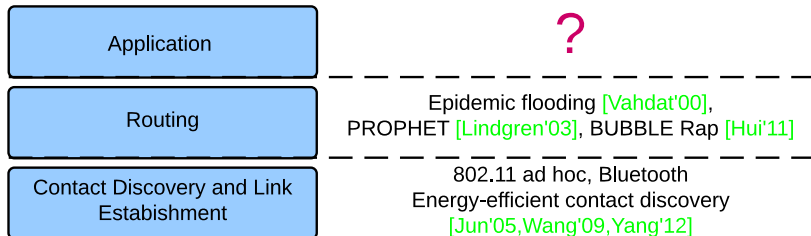


## Characteristics

- No end-to-end connectivity & long delay
- Resource-limited network nodes



# DTNs Related Work



# Our Contributions: Deployment-Based DTN Performance Characterization

## Utility

Message delivery performance characterization

## Robustness

Performance evaluation under blocking and censorship attacks

## Overhead

Energy overhead modeling and evaluation

- measurement-based

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## Why A New Deployment?

Existing traces are outdated...

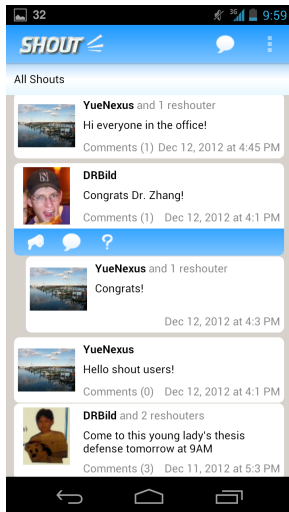
	UCSD	Dartmouth
Year	2002	2003/2004
Device	PDA	Laptop/PDA
Duration (days)	77	114
# of devices	273	6,648

Likely **cannot** capture the mobility of modern smartphone users

- Dartmouth: mostly laptops
- UCSD: limited-functionality PDAs

User engagement declined during study

# 1am: A Microblogging Service on DTN

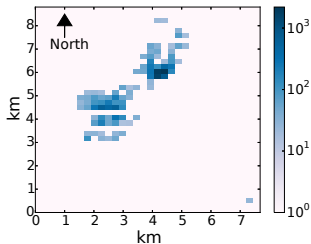
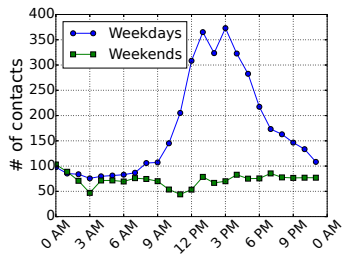
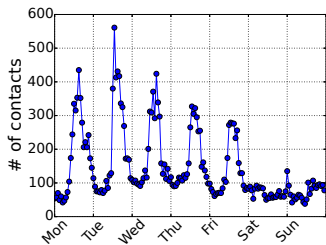


- A Twitter-like microblogging app
- Installed by UMich students and faculty members (291 in total)
- WiFi based contact estimation

# 1am Trace Summary

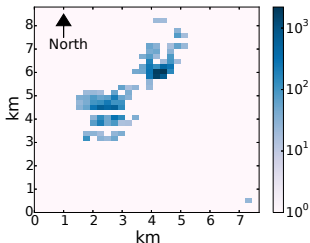
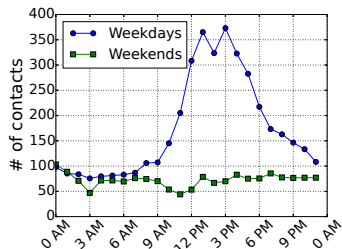
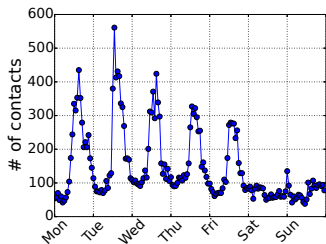
	UCSD	Dartmouth	1am
Year	2002	2003/2004	2013
Device	PDA	Laptop/PDA	Smartphone
Duration (days)	77	114	31
Granularity (seconds)	120	300	114.5
Devices participating	273	6,648	111
# contacts/pair/hour	0.0028	0.000067	0.033

# Temporal and Spatial Patterns of Contacts





# Temporal and Spatial Patterns of Contacts



- Work- or study- related co-location
- Concentrated in very few buildings in the north campus

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# Evaluation Methodology: Trace-Based Simulation

## Why simulation?

- Not enough user initiated messages
- Based on users' contact traces

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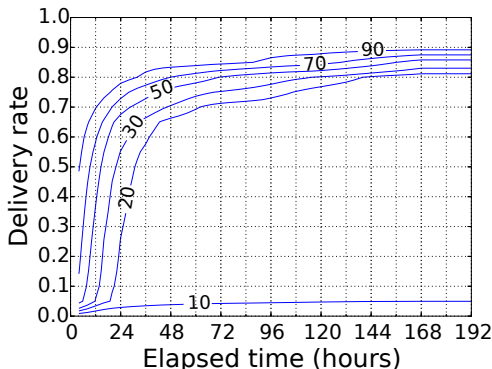
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## Setup

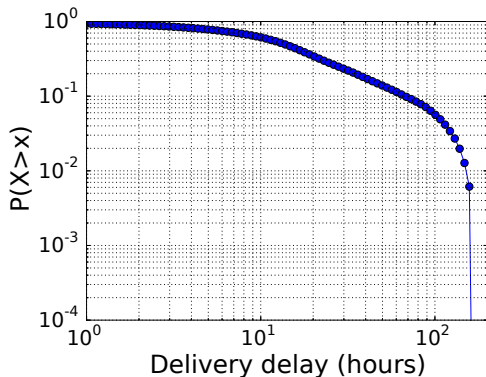
- 1 message per device per 6 minute
- Epidemic flooding

# Message Delivery Progress



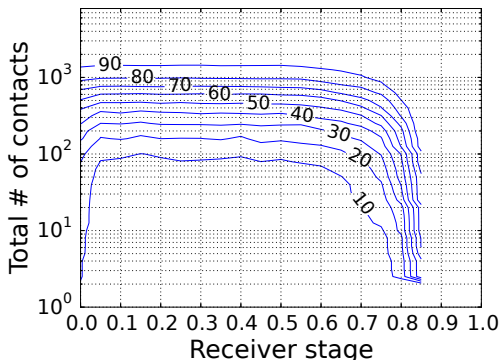
- Median delivery rate: 0.85
- Fast to reach first 60%-70% of nodes
- Extremely hard to reach the rest ?Why?

# Delivery Delay Distribution



- Power-law distribution
- Median delay: 13 hours

## Why So Hard to Reach The Last 30%?



- Contact heterogeneity causes the change in delivery speed
- The last 30% have diminishing contacts

## Summary on Performance

111 users, 4 km $\times$ 4 km campus (only 0.2% adoption rate)

- Median delivery rate after a day: 0.68
- Median delivery rate after a week: 0.85
- Median delay: 13 hours
- Extreme long delay for the last 20%-30% nodes



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### Significant performance variations

- Significant performance variations caused by contact heterogeneity
- **ATTN: Mobility models resulting in contact homogeneity are insufficient!**

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# How the Network Reacts to Resource Removal Attacks?

## Resource removal

- Blocking attacks = removing network resources
- Censorship = a selective form of blocking attacks

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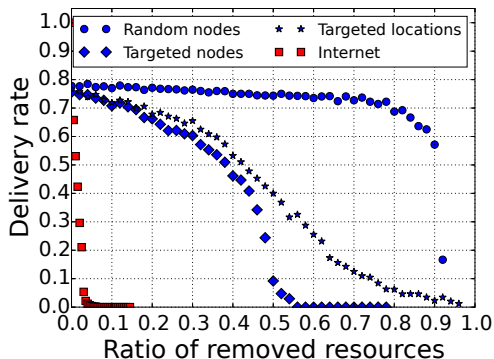
## Resource removal

- Blocking attacks = removing network resources
- Censorship = a selective form of blocking attacks

## Attack strategies

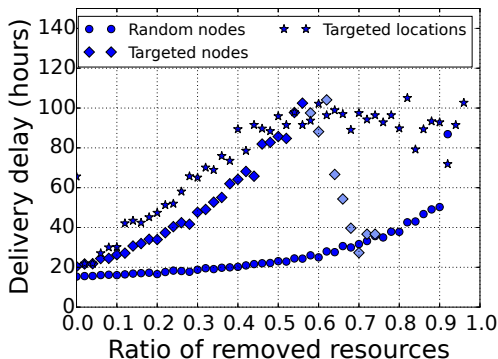
- **Random device removal:** Remove devices at random
- **Targeted device removal:** Remove devices appearing most frequently on shortest paths
- **Targeted location removal:** Remove locations appearing most frequently on shortest paths

# Message Delivery Rate Degradation Under Attacks



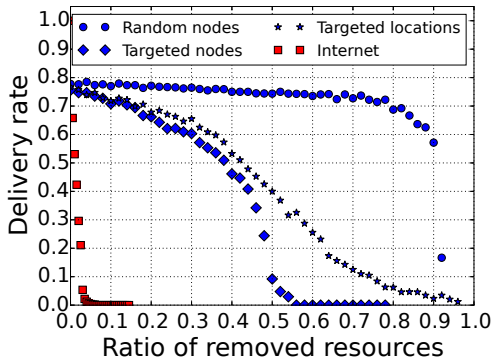
- Random device removal: almost no effect until more than 80% removed
- Targeted removal: fast degradation after more than 30%-40% removed
- The Internet: quick fragmentation after 3% removed [Albert'00]

# Message Delivery Delay Degradation Under Attacks



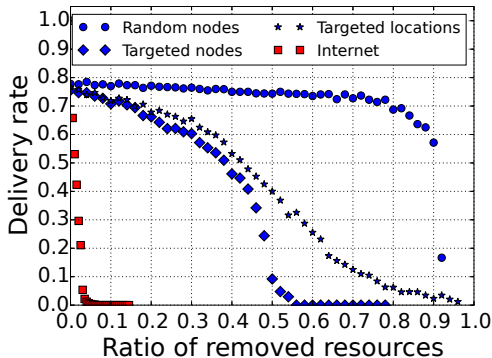
- Random device removal: almost no effect until more than 60% removed
- Targeted removal: faster, gradual degradation
- Shaded points: dropped messages are not included

## Summary on Robustness



- Non-hierarchical networks are much more robust
- Caveat: high-hierarchy Internet routers may be hard to attack...

## Summary on Robustness



- Non-hierarchical networks are much more robust
- Caveat: high-hierarchy Internet routers may be hard to attack...
- Question: Will the same properties hold when the network scales up, e.g. with 1,000, or 10,000 users?



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# Energy Model of All-to-All Content Sharing

## Contact discovery

- Constant

## Communication\*

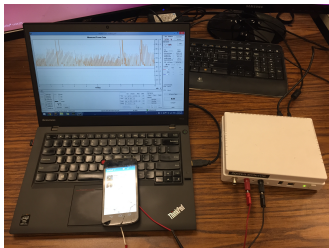
- Message delivery:  $\mathcal{O}(N)$
- Metadata exchange:  $\mathcal{O}(N^2)$

## Computation

- $\mathcal{O}(N)$

\*The epidemic flooding protocol is used

# Energy Model of the Wireless Component in 802.11 Ad Hoc Mode



$P_{idle}$ (idle state)	210 mW
$P_{high}$ (send/receive state)	341 mW
$E_{sw}$ (one power state switch from low-high-low)	242 mJ
$E_{ping}$ (send+receive one 802.11b frame (ping))	2.16 mJ
$E_{send}$ (send one 802.11b frame (UDP))	19.6 mJ
$E_{recv}$ (receive one 802.11b frame (UDP))	6.8 mJ

## Different Application Message Types

### Message Sizes ( $F_m$ ) and Initiation Frequencies ( $f_m$ )

Type	Typical Message size	$F_m$ (frames/msg.)	$f_m$ (msg./hour)
Text	514 characters	1	0.16667*
Image	100 KB	44	0.02536**
Video clip	1 Mbps $\times$ 240 seconds	12000	0.00054***

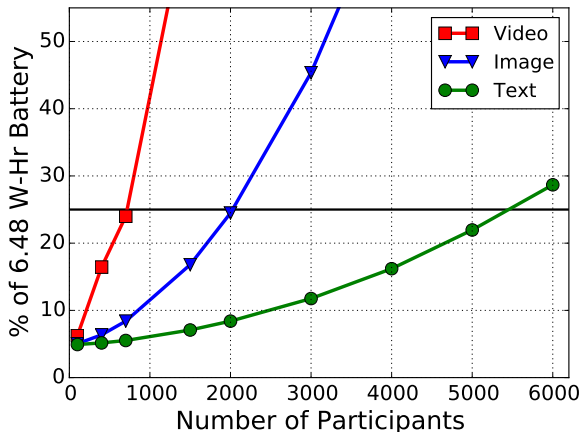
\*This is Twitter's average rate

\*\*This is Facebook's average photo posting rate

\*\*\*This is YouTube's average rate

## Results

### 12 Hour Energy Consumption against Network Size



• Text & Image: 1,500 nodes

• All Three: 400 nodes

# Discussions on Energy Overhead

## Battery energy severely limits network scale

- If everyone in our university joins to text, battery runs out in **less than 1 hour!**
- If everyone in our university joins to send photos, battery runs out in **less than 10 minutes!**

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## Discussions

- Dominating contributor: **metadata exchange**
- How to reduce the cost of metadata exchange?

## Conclusion and Future Work

Promising performance with only 0.2% adoption rate

- ~70% delivery rate within 24 hours
- Median delay 13 hours



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- The network breaks down after more than 40% are compromised

## Conclusion and Future Work

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- Median delay 13 hours

### Robustness against paralyzing and censorship attacks

- The network breaks down after more than 40% are compromised

### Battery energy severely limits scale

- Cannot support the entire university population